Learning Concepts Described by Weight Aggregation Logic

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Learning from Examples





Input: Labelled Examples

labelled tuples (v_1, \ldots, v_k)



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labelled tuples (v_1, \ldots, v_k)

Output: Concept

formula $\varphi(x_1, \ldots, x_k; y_1, \ldots, y_\ell)$, parameters w_1, \ldots, w_ℓ





Positive Examples Bob Emma

Negative Examples



Positive Examples Bob Emma

Negative Examples

Possible Concept

Carol's friends



Positive Examples Bob Emma

Negative Examples Dan

Possible Concept

Carol's friends

$$\varphi(x;y) = F(x,y)$$

parameter: Carol



Positive Examples (Alice, Bob) (Dan, Emma) (Carol, Dan)



Positive Examples (Alice, Bob) (Dan, Emma) (Carol, Dan)

2

Negative Examples

(Bob, Dan) (Alice, Carol)



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Negative Examples (Bob, Dan)

(Alice, Carol)

Possible Concept

having a common friend who is not Emma



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$$\varphi(x_1, x_2; y) = \exists z \left(F(x_1, z) \land F(x_2, z) \land z \neq y \right)$$

parameter: Emma

Theorem (Grohe and Ritzert, 2017)

Concepts definable in FO on structures of small degree can be learned in sublinear time (in the size of the structure).

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Concepts definable in the weight aggregation logic FOWA₁ on weighted structures of small degree can be learned in sublinear time with quasilinear precomputation.

Learning Concepts in Sublinear Time

We only consider formulas of limited complexity (limited nesting depth)

- 1: for all formulas do
- 2: for all parameters do
- 3: **if** concept is consistent **then**
- 4: **return** concept
- 5: **reject**

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- 1) Gaifman normal form, uses enriched structure with quasilinear-time precomputation

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small degree, small neighbourhood (polylogarithmic)

- 1) Gaifman normal form
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Weight Aggregation Logic



- relational structure
- + weights for weight symbols in ${f W}$



- relational structure
- + weights for weight symbols in W

Example

- $W = \{a, \ell\}$
- age $a(v) \in \mathbb{Z}$
- length of friendship $\ell(v, w) \in \mathbb{Z}$



Terms

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$$t_1(x, y) = 5 \cdot a(x) - a(y) \cdot \ell(x, y)$$



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 $FOWA_1$



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FOWA₁

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FOWA₁

- only finite ring in $s = \sum w(y_1, \dots, y_k).\varphi$
- only one free variable in $P(t_1, \ldots, t_m)$

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Localisation Theorem for $\ensuremath{\mathsf{FOWA}}\xspace_1$

Use local formula from FOW_1 on enriched structure instead.

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Gaifman normal form for $\ensuremath{\mathsf{FOW}}_1$

Feferman-Vaught decompositions for FOW_1

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Future work

allow more free variables in $P(t_1, \ldots, t_m)$ in a guarded setting