



# Learning Aggregate Queries Defined by First-Order Logic with Counting

Steffen van Bergerem and Nicole Schweikardt

ICDT 2025

# How to be a Good Colleague

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Name	Popularity
Alice	5
Bob	1
Carol	2
Dan	3
Emma	1

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(names changed for privacy reasons)

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Name	Popularity
Alice	5
Bob	1
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Name	Type of Cake
Alice	chocolate
Dan	lemon
Carol	strawberry
Alice	chocolate
Bob	carrot
Emma	apple
Dan	chocolate
Alice	strawberry
Carol	lemon

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# How to be a Good Colleague

Name	Popularity
Alice	5
Bob	1
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Emma	1

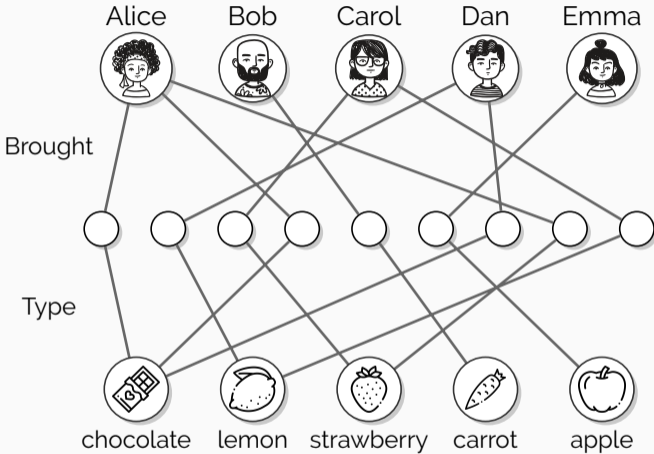
(names changed for privacy reasons)

Name	Type of Cake
Alice	chocolate
Dan	lemon
Carol	strawberry
Alice	chocolate
Bob	carrot
Emma	apple
Dan	chocolate
Alice	strawberry
Carol	lemon

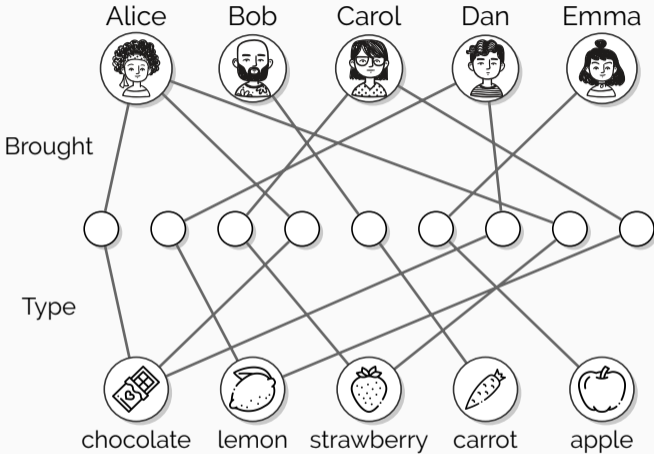
$$\text{Popularity} = 2 \cdot \#\text{chocolate cakes} + \#\text{other cakes}$$

# How to be a Good Colleague

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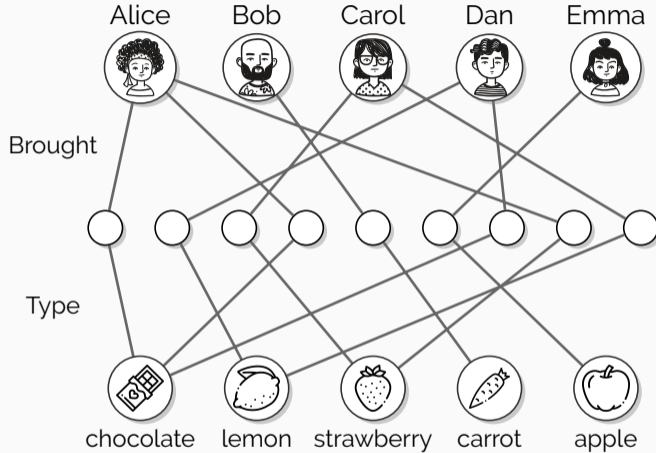
# How to be a Good Colleague



- (Alice, 5)
- (Bob, 1)
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- (Emma, 1)



# How to be a Good Colleague



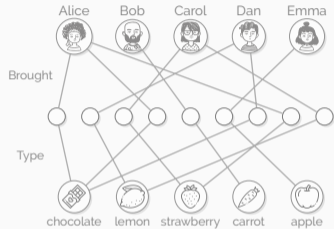
- (Alice, 5)
- (Bob, 1)
- (Carol, 2)
- (Dan, 3)
- (Emma, 1)

$$p(x) = 2 \cdot \#(c).(\text{Brought}(x, c) \wedge \text{Type}(c, \text{chocolate})) + \#(c).(\text{Brought}(x, c) \wedge \neg \text{Type}(c, \text{chocolate}))$$

# Learning from Examples

## Precomputation phase

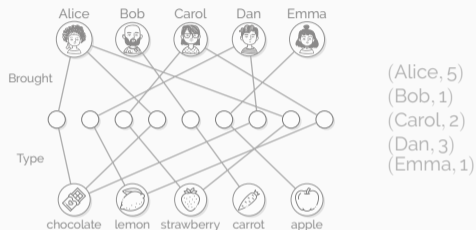
Given a database  $D$ , build index structure



# Learning from Examples

## Precomputation phase

Given a database  $D$ , build index structure



## Learning phase

Given list of labelled examples  $(\bar{v}, \lambda) \in (U(D))^k \times \mathbb{Z}$

Return term  $t(\bar{x}) \in \text{FOC}_1$  (of certain maximum complexity)  
such that  $\llbracket t(\bar{v}) \rrbracket^D = \lambda$  for all given examples  $(\bar{v}, \lambda)$

or reject if there is no such term

## Results on databases of polylogarithmic degree

### Grohe and Ritzert, LICS 2017

Boolean-valued concepts definable in [first-order logic](#) can be learned in sublinear time.

### v. B. and Schweikardt, CSL 2021

Boolean-valued concepts definable in [first-order logic with counting](#) or [first-order logic with weight aggregation](#) can be learned in sublinear time after quasi-linear-time precomputation.

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### v. B. and Schweikardt, ICDT 2025

[Integer-valued](#) concepts definable in [first-order logic with counting](#) can be learned in sublinear time after quasi-linear-time precomputation.

Main tool: locality results similar to Gaifman normal forms

# First-Order Logic with Counting (FOC)

Counting terms

Formulas

# First-Order Logic with Counting (FOC)

## Counting terms

- $t_{\exists c}(x) = \#(c).(Brought(x, c) \wedge Type(c, \exists c))$
- $t_{\neg \exists c}(x) = \#(c).(Brought(x, c) \wedge \neg Type(c, \exists c))$

## Formulas

# First-Order Logic with Counting (FOC)

## Counting terms

- $t_{\exists}(x) = \#(c).(Brought(x, c) \wedge Type(c, \exists))$
- $t_{\neg\exists}(x) = \#(c).(Brought(x, c) \wedge \neg Type(c, \exists))$
- $p(x) = 2 \cdot t_{\exists}(x) + t_{\neg\exists}(x)$

## Formulas



# First-Order Logic with Counting (FOC)

## Counting terms

- $t_{\text{obj}}(x) = \#(c).(Brought(x, c) \wedge Type(c, \text{obj}))$
- $t_{\neg \text{obj}}(x) = \#(c).(Brought(x, c) \wedge \neg Type(c, \text{obj}))$
- $p(x) = 2 \cdot t_{\text{obj}}(x) + t_{\neg \text{obj}}(x)$

## Formulas

- $\varphi_1(x) = t_{\text{obj}}(x) \leq t_{\neg \text{obj}}(x)$
- $\varphi_2(x) = \forall y (t_{\text{obj}}(y) \leq t_{\text{obj}}(x))$

# First-Order Logic with Counting (FOC)

## Counting terms

- $t_{\text{type}}(x) = \#(c).(Brought(x, c) \wedge Type(c, \text{type}))$
- $t_{\neg \text{type}}(x) = \#(c).(Brought(x, c) \wedge \neg Type(c, \text{type}))$
- $p(x) = 2 \cdot t_{\text{type}}(x) + t_{\neg \text{type}}(x)$
- $q = \#(x) \cdot \varphi_1(x)$

## Formulas

- $\varphi_1(x) = t_{\text{type}}(x) \leq t_{\neg \text{type}}(x)$
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# First-Order Logic with Counting (FOC)

## Counting terms

- $t_{\text{obj}}(x) = \#(c).(Brought(x, c) \wedge Type(c, \text{obj}))$
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- $p(x) = 2 \cdot t_{\text{obj}}(x) + t_{\neg \text{obj}}(x)$
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## Formulas

- $\varphi_1(x) = t_{\text{obj}}(x) \leq t_{\neg \text{obj}}(x)$
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## FOC<sub>1</sub>

- introduced by Grohe and Schweikardt (PODS 2018)
- subformulas comparing terms have at most one free variable
- has Gaifman-style normal forms

# Main result

## v. B. and Schweikardt, ICDT 2025

For **integer-valued** concepts definable in the **first-order logic with counting**  $\text{FOC}_1$ , there is an algorithm for the learning problem with

- **precomputation phase** in  $n \cdot d^{\mathcal{O}(1)}$
- **learning phase** in  $(d + t)^{\mathcal{O}(1)}$

( $n$ : size of active domain,  $d$ : degree of the database,  $t$ : number of examples)

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For databases of polylogarithmic degree:

- **precomputation phase** in quasi-linear time  $n \cdot (\log n)^{\mathcal{O}(1)}$
- **learning phase** in time polylogarithmic in the size of the database  $(\log n + t)^{\mathcal{O}(1)}$

## Discussion

v. B. and Schweikardt, ICDT 2025

Integer-valued concepts definable in first-order logic with counting can be learned in polylog time after quasi-linear-time precomputation.

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Integer-valued concepts definable in first-order logic with counting can be learned in polylog time after quasi-linear-time precomputation.

- data complexity vs. parameterised complexity
- consistent learning vs. probably approximately correct (PAC) learning
- logics with weight aggregation

**Bring more (chocolate) cakes!**

